

Scalable Abstractions for Scalable Systems

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Socialite

Parallel program = Operator + Schedule + Parallel data structure **OPERATOR FORMULATION NATIVE PERFORMANCE**

Active element — Site where computation is



needed

Operator — Computation at active element

Activity

— application of operator to active element

active node neighborhood

Neighborhood

- Set of nodes/edges read/written by activity
- Distinct usually from neighbors in graph

Ordering

- Scheduling constraints on execution order of activities
- Unordered algorithms: no semantic constraints but performance may depend on schedule
- Ordered algorithms: problem-dependent order

Amorphous data-parallelism



🛯 Combblas 🖾 Graphlab

🛛 Giraph

480

448

384

352

320

288

256

224

192

160

128

delaunay mesh

Galois 🏼

code.

"Navigating the maze of graph analytics frameworks" Nadathur et al SIGMOD 2014

Intel HPC Study: Galois

implementations are

comperable to hand-

written and optimized

Combblas 🛛 Graphlab

Galois 🎆

edia

barnes-hut launay mesh refinement delaunay triangulation betweenness centrality triangles	Numa-awareruntime and datastructures forhigh multi-corescaling. Realperformancesingle-threadperformance isperformance ison par withoptimized serial code
<u></u>	optimized serial code
64 128 192 256 320 384 448 512	
Inreads	

App	Implementation	Threads	Time (s)
dmr	triangle	1	96
	Galois	1	155.7
	Galois	512	0.37
dt	triangle	1	1185
	Galois	1	56.6
	Galois	512	0.18
bh	splash2	1	>6000
	Galois	1	1386
	Galois	512	3.55
bc	HPCS SSCA	1	6720
	Galois	1	5394
	Galois	512	21.6
tri	graphlab	2	531
	Galois	1	7.03
	Galois	512	0.028

ble 2: Serial runtime comparisons to other imple itations rounded to the nearest second. Included runtimes for Galois algorithms at 512 threads e splash2 implementation of bh timed out after 0 minutes

THINK BEYOND A VERTEX

— Multiple active nodes can be processed in parallel subject to neighborhood and ordering constraints

MULTI-LEVEL PROGRAMMING

Joe Programmer: **Operator and Schedule** Specification

Stephanie Programmer: Parallel Data Structures

Obi-Wan Programmer: Synchronization, NUMA, Scalable Runtime



The best algorithm may not be expressible as a vertex program

- Connected components with union-find
- The best algorithm may require application-specific scheduling — Priority scheduling for SSSP

Autonomous scheduling required for highroad diameter graphs

— Coordinated scheduling uses many rounds and has too much overhead



twitter50



ONE MODEL, MANY TARGETS

GPUs

- Optimized implementation strategies for coalesing, synchronization, high-thread counts
- Scheduling and load balancing for highly-parallelism

Ubiquitous parallelism:

— small number of expert programmers (Stephanies) must support large number of application programmers (Joes) — cf. SQL

Stephanie

- Library of concurrent data structures
- Provides serializable, atomic execution of activities

Joe

- Application code in stylized, sequential C++
- Uses Galois set iterator for highlighting opportunities for exploiting ADP and Galois data structures for concurency control

— Multi-gpu support

FPGAs

- Optimized scheduling and conflict detection
- Targets network of FPGAs architectures

Distributed Memory

— Transparent support for distributed memory for arbitrarily complex irregular algorithms

Xeon PHI

— Numa and memory optimizations lead to out-performing simple pthread/openMP/openCL codes

Heterogeneous and emerging — Mixed GPU/CPU

- In progress: Cluster of multicore/CPU, coherent CPU/FPGA